

FRIDAY TALKS @ TUE.AI #300126

THE VALUE OF AMBIGUOUS COMMITMENT WITH MULTIPLE FOLLOWERS

THIS IS JOINT WORK WITH...



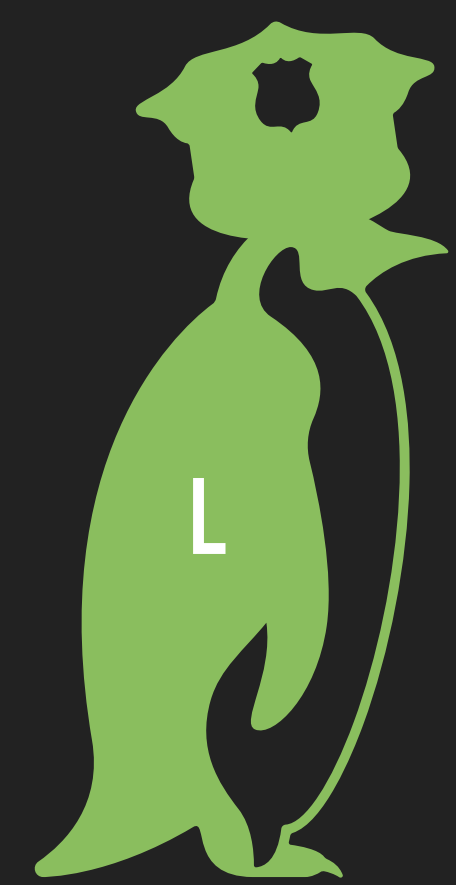
Natalie Collina



Aaron Roth



SIMPLE STACKELBERG GAME



$$\mathcal{A}_L := \{a, b, c\}$$

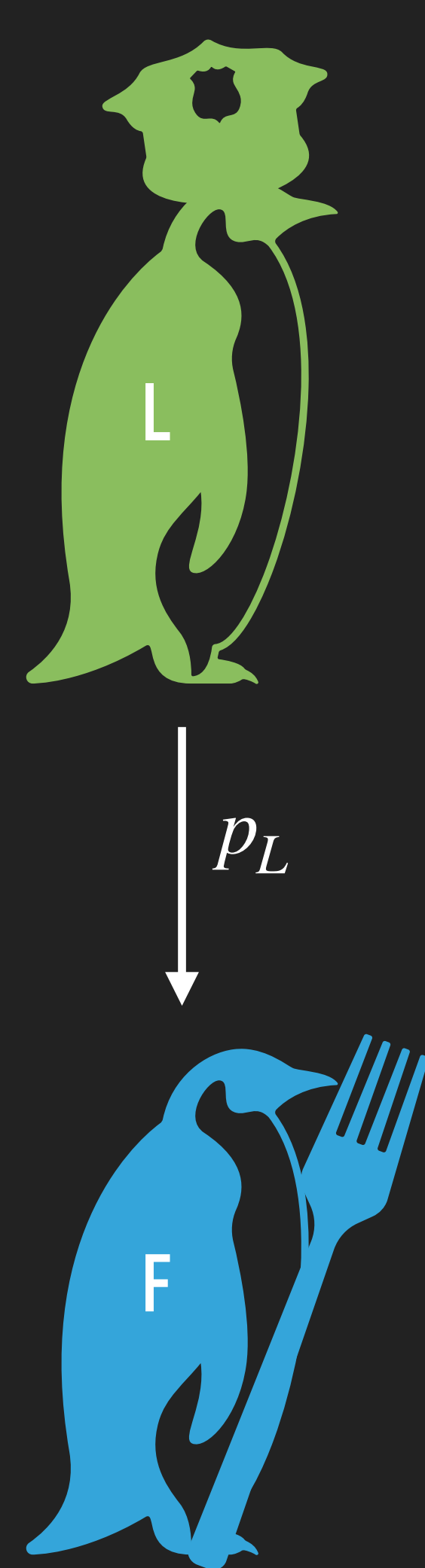
L	a	b	c
i	$u(a,i)$	$u(b,i)$	$u(c,i)$
ii	$u(a,ii)$	$u(b,ii)$	$u(c,ii)$



$$\mathcal{A}_F := \{i, ii\}$$

F	a	b	c
i	$v(a,i)$	$v(b,i)$	$v(c,i)$
ii	$v(a,ii)$	$v(b,ii)$	$v(c,ii)$

SIMPLE STACKELBERG GAME



L	a	b	c
i	$u(a,i)$	$u(b,i)$	$u(c,i)$
ii	$u(a,ii)$	$u(b,ii)$	$u(c,ii)$

F	a	b	c
i	$v(a,i)$	$v(b,i)$	$v(c,i)$
ii	$v(a,ii)$	$v(b,ii)$	$v(c,ii)$

1. Leader chooses $p_L \in \Delta(\mathcal{A}_L)$

2. Follower best responds

$$\text{BR}(p_L) = \arg \max_{a \in \mathcal{A}_F} V(p_L, a) *$$

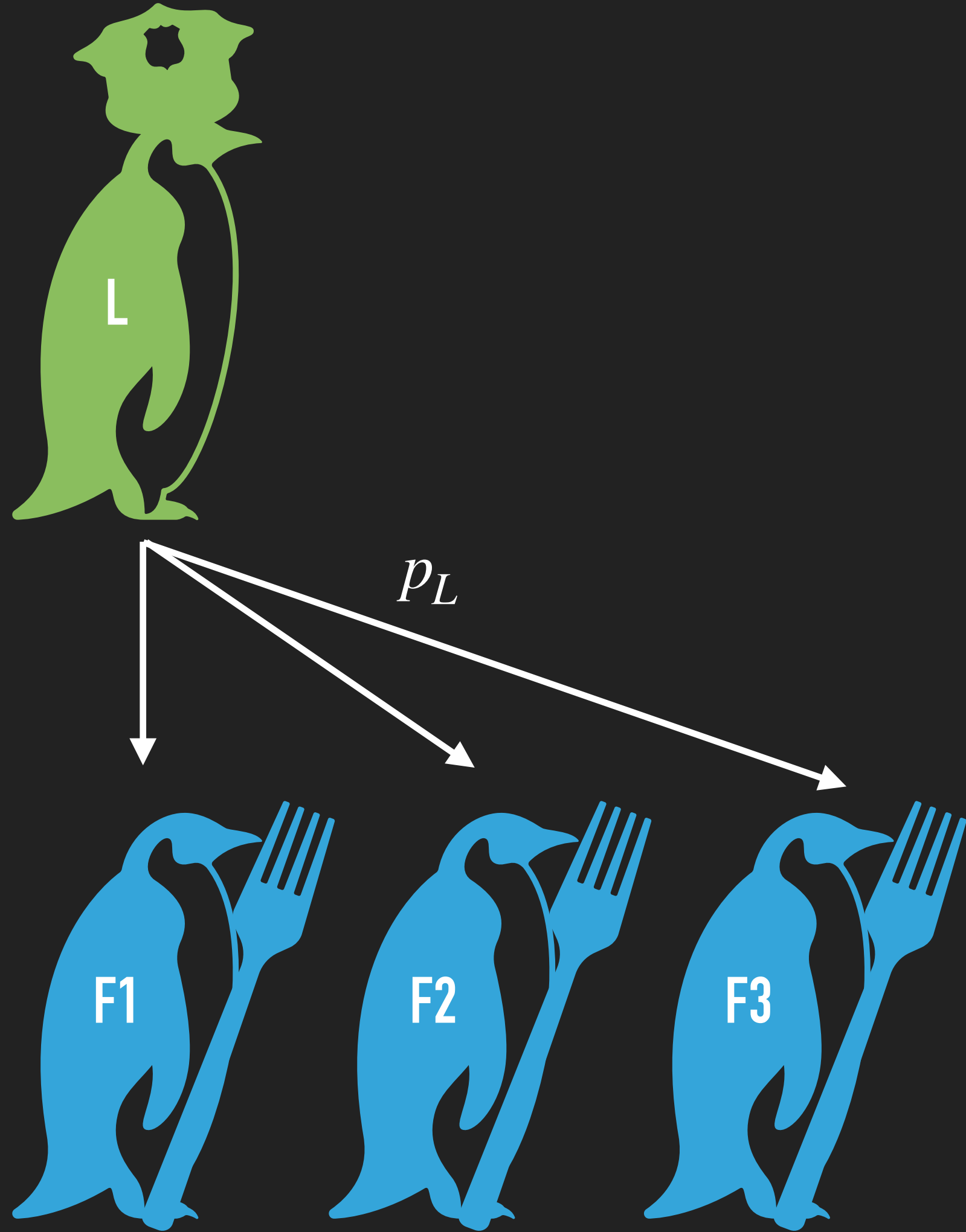
3. Leader obtains the Stackelberg value $U(p_L, \text{BR}(p_L))^{**}$

* $V(p, q) = \mathbb{E}_{a_L \sim p, a_F \sim q}[v(a_L, a_F)]$

** $U(p, q) = \mathbb{E}_{a_L \sim p, a_F \sim q}[u(a_L, a_F)]$

COUPLED STACKELBERG GAME

$$\forall F \in \mathcal{F}, u_F: \mathcal{A}_L \times \mathcal{A}_F \rightarrow \mathbb{R}$$



$$\forall F \in \mathcal{F}, \quad v_F: \mathcal{A}_L \times \mathcal{A}_F \rightarrow \mathbb{R}$$

1. Leader chooses $p_L \in \Delta(\mathcal{A}_L)$

2. **All** followers best respond

$$\text{BR}_F(p_L) = \arg \max_{a \in \mathcal{A}_F} V_F(p_L, a), \forall F \in \mathcal{F}$$

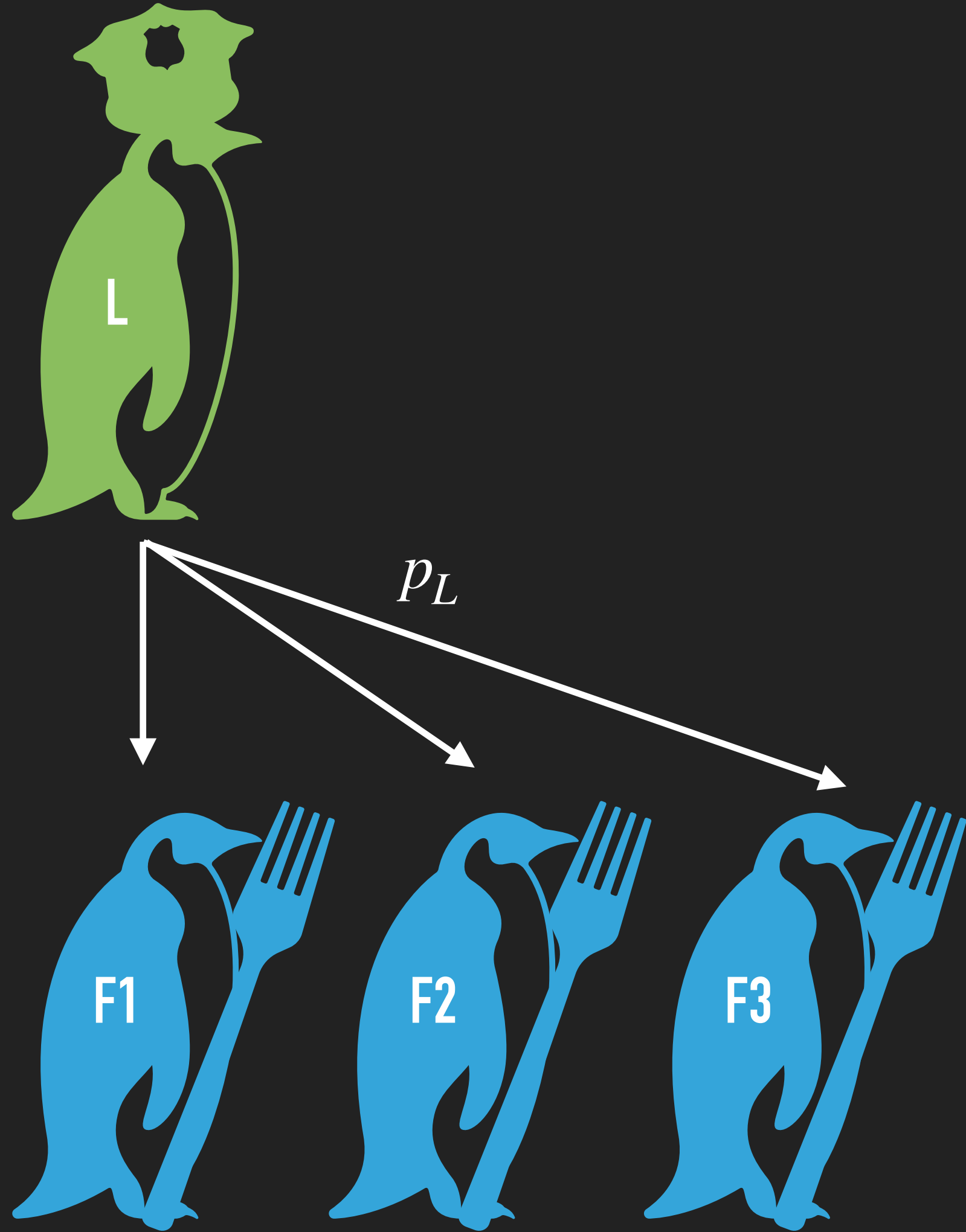
3. Leader obtains the Stackelberg value

$$\sum_{F \in \mathcal{F}} U_F(p_L, \text{BR}_F(p_L))$$

**AMBIGUOUS
COMMITMENT**

COUPLED STACKELBERG GAME WITH AMBIGUOUS COMMITMENT

$$\forall F \in \mathcal{F}, u_F: \mathcal{A}_L \times \mathcal{A}_F \rightarrow \mathbb{R}$$



$$\forall F \in \mathcal{F}, v_F: \mathcal{A}_L \times \mathcal{A}_F \rightarrow \mathbb{R}$$

1. Leader chooses $P_L \subseteq \Delta(\mathcal{A}_L)$, closed cvx.

2. All followers best respond

$$\text{BR}_F(P_L) = \arg \max_{p \in \Delta(\mathcal{A}_F)} \min_{p_L \in P_L} V_F(p_L, a), \forall F \in \mathcal{F}$$

3. Leader obtains the Stackelberg value

$$\min_{p_L \in P_L} \sum_{F \in \mathcal{F}} U_F(p_L, \text{BR}_F(P_L))$$

THEOREM I (INFORMAL)

If there is only a single follower, the best classical commitment achieves an equal or higher Stackelberg value than any ambiguous commitment does.

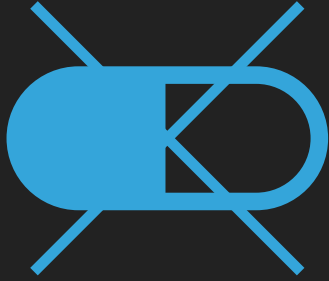
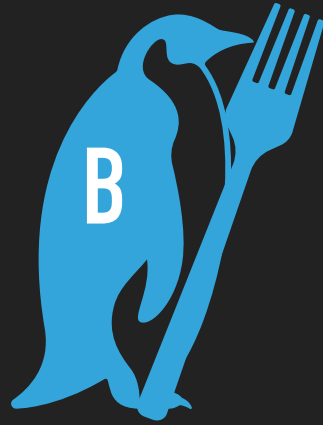
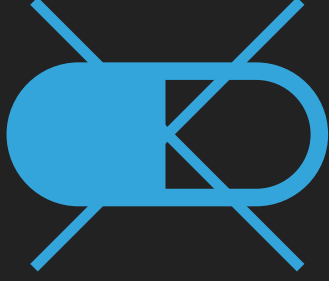
BUT

How about a coupled Stackelberg game with multiple followers?

A FIRST FARM CONTROL GAME



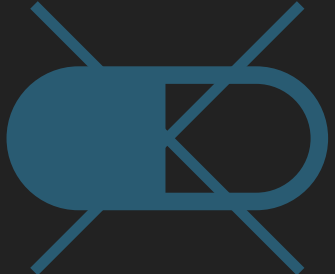
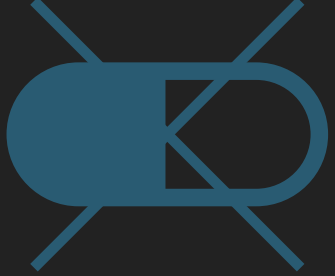
Zero-sum game with
two followers
(followers maximize,
leader minimizes)

	 			
	1	1	1	1
	0	0	4	4
	 			
	1	1	1	1
	0	4	0	4

A FIRST FARM CONTROL GAME



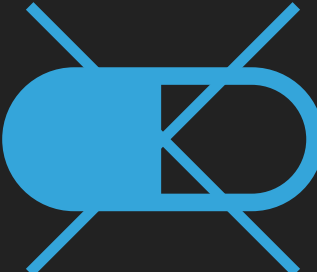
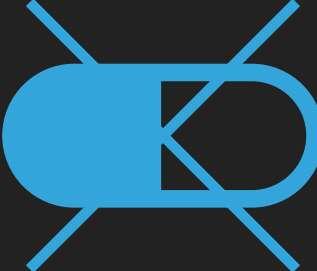
Zero-sum game with
two followers
(followers maximize,
leader minimizes)

	 			
	1	1	1	1
	0	0	4	4
RESOURCE-CONSTRAINT				
	 			
	1	1	1	1
	0	4	0	4

A SECOND FARM CONTROL GAME



Zero-sum game with
two followers
(followers maximize,
leader minimizes)

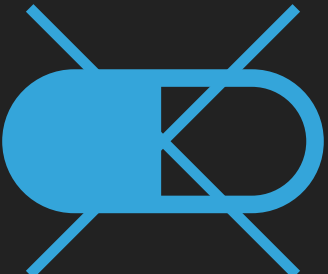
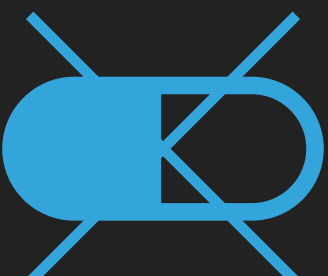
		
	1	1
	4	0
		
	1	1
	0	4

A SECOND FARM CONTROL GAME

EVOCATION OF RESPONSE PATTERN UNREACHABLE
BY CLASSICAL COMMITMENT



Zero-sum game with
two followers
(followers maximize,
leader minimizes)

		
	1	1
	4	0
		
	1	1
	0	4

THEOREM II (INFORMAL)

There exists a two follower zero-sum game in which ambiguous commitment leads to an arbitrarily higher Stackelberg value than any classical commitment.

OUR COMPUTATIONAL RESULTS

1. A poly-runtime algorithm for finding the optimal ambiguous commitment in a $2 \times N$ -Stackelberg games with M followers.
2. NP-Hardness of finding the optimal ambiguous commitment in $K \times N$ -Stackelberg games with a fixed number of followers.

**“AMBIGUOUS COMMITMENTS HELP TO
RAISE THE STACKELBERG VALUE IN
MULTIPLE FOLLOWER GAMES.”**

— WINSTON CHURCHILL . . . OR MAYBE SOMEONE ELSE